Distributed Agent-Based Ant Colony Optimization for Solving Traveling Salesman Problem on a Partitioned Map

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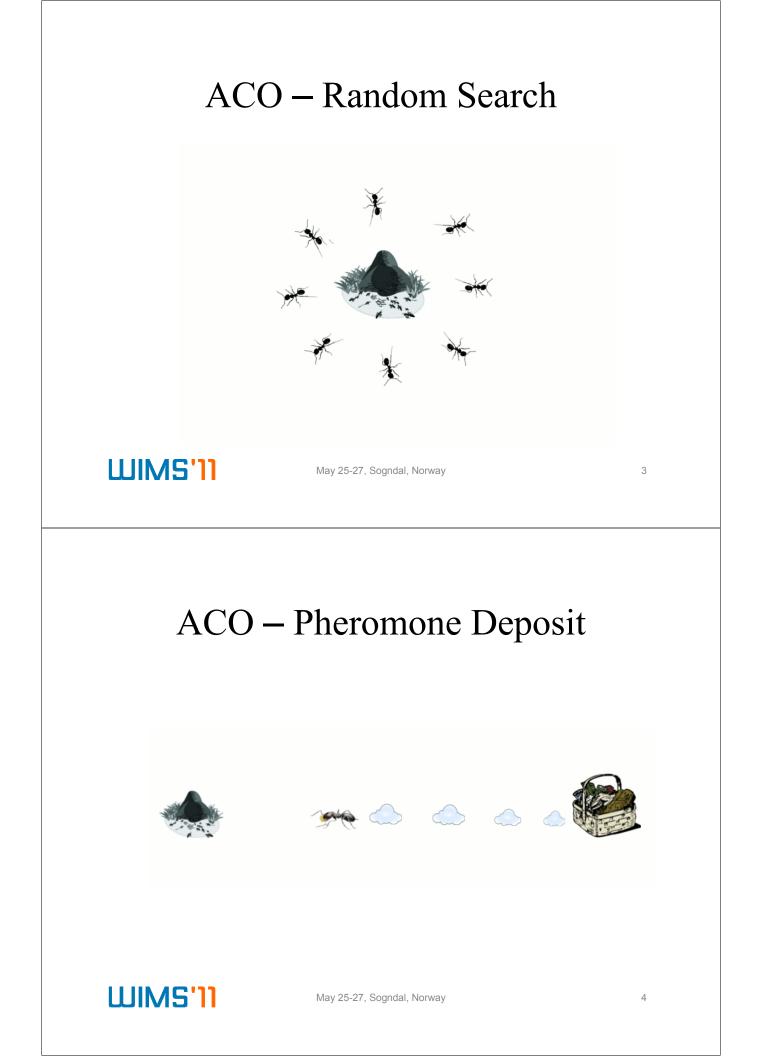
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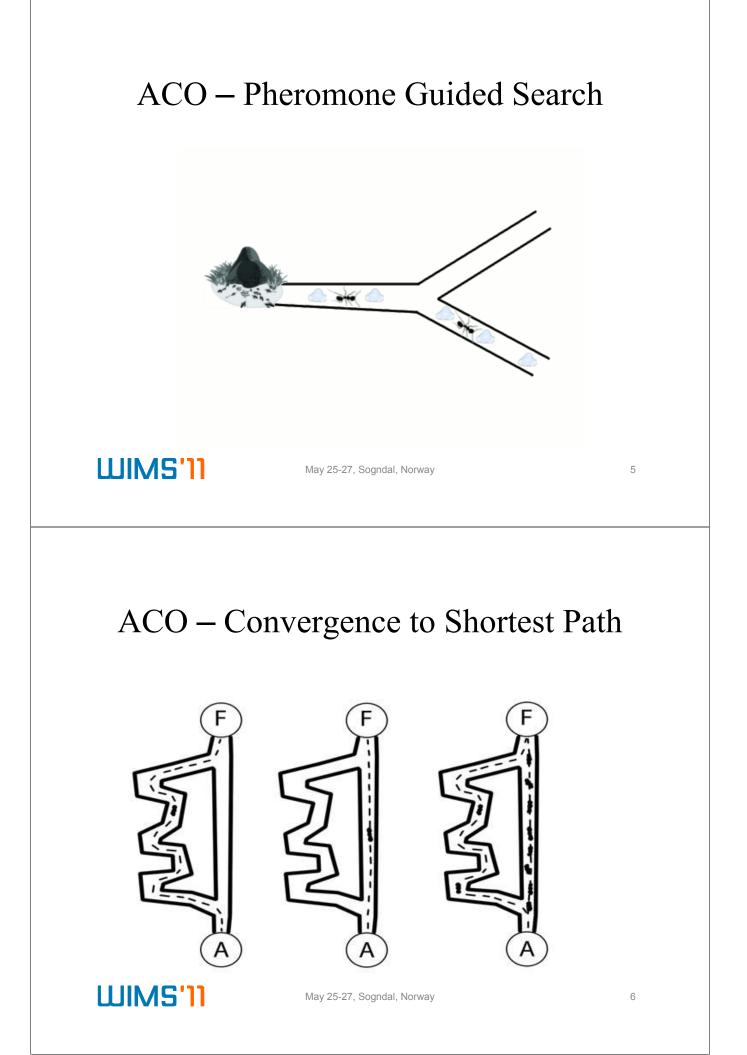
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Overview

- Ant Colony Optimization
- Mathematical Model
- Distributed Architecture => ACODA
- The Traveling Salesman Problem
- Experimental results

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Approaches for distributing ACO

| | Sequential | Fine-grained | Multi-colony | Agent-based |
|--------------|---------------|--|--|--|
| Environment | data | data | data | a set of agents |
| Ants | functions | functions | functions | a set of agents (sometimes mobile) |
| Messages | none | ants synchronize after each iteration | colonies synchronize with a given frequency | ants communicate with each other and with nodes |
| Ant Movement | one at a time | all at once synchronized | asynchronous (but one at a time in each colony) | asynchronous |
| Observations | | large overhead | currently best approach | large overhead |

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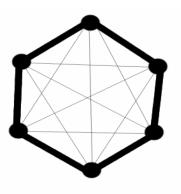
Motivation

- Use of multi-agent systems for modeling ants' environment.
 - It was observed that complexity of ants' movement stems from the complexity of the environment.
- Mapping of ants' environment to a distributed architecture and the mapping of the ants' migration to messages exchanged between the agents located in the ants' environment.
 - *n* agents, *N* ant migrations/any 2 agents, cost of ant message on a single machine = a and between 2 machines = b > a.
 - Execution time on 1 machine $T_1 = a N n (n-1)/2$ and on n machines $T_n = b N (n-1)$. If $n \ge 2b/a$ then $T_1 \ge T_n$.



Traveling Salesman Problem

• Given a weighted graph, the goal is to find the shortest tour that visits each node exactly once.



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Probabilistic Choices

$$p_{i,j} = \frac{(\tau_{i,j}^{\alpha})(\eta_{i,j}^{\beta})}{\sum (\tau_{i,j}^{\alpha})(\eta_{i,j}^{\beta})}$$

where:

- $\tau_{i,j}$ = amount of pheromone deposited on edge (*i*,*j*)
- $\alpha =$ parameter to control the influence of $\tau_{i,j}$
- $\eta_{i,j} = desirability$ of edge (i,j) computed as the inverse of the weight $w_{i,j}$ of edge (i,j), i.e.
- β = parameter to control the influence of $\eta_{i,j}$
- *j* = a node reachable from node *i* that was not visited yet

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Pheromone Increment

 $\Delta \tau_{i,j}^{k} = \begin{cases} 1/L_{k} \text{ if ant } k \text{ travels on edge}(i,j) \\ 0 \text{ otherwise} \end{cases}$

where:

- L_k is the cost of the *k*-th ant tour.
- Δτ is the amount of pheromone ant k deposits on edge (i,j)

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Pheromone Deposit

$$\tau_{i,j} = (1 - \rho)\tau_{i,j} + \rho \Delta \tau_{i,j}^k$$

where:

- $\tau_{i,j}$ is the amount of pheromone on edge (i,j)
- Δτ is the amount of pheromone ant k deposits on edge (i, j)
- ρ is the evaporation rate $0 \le \rho < 1$

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Local Evaporation

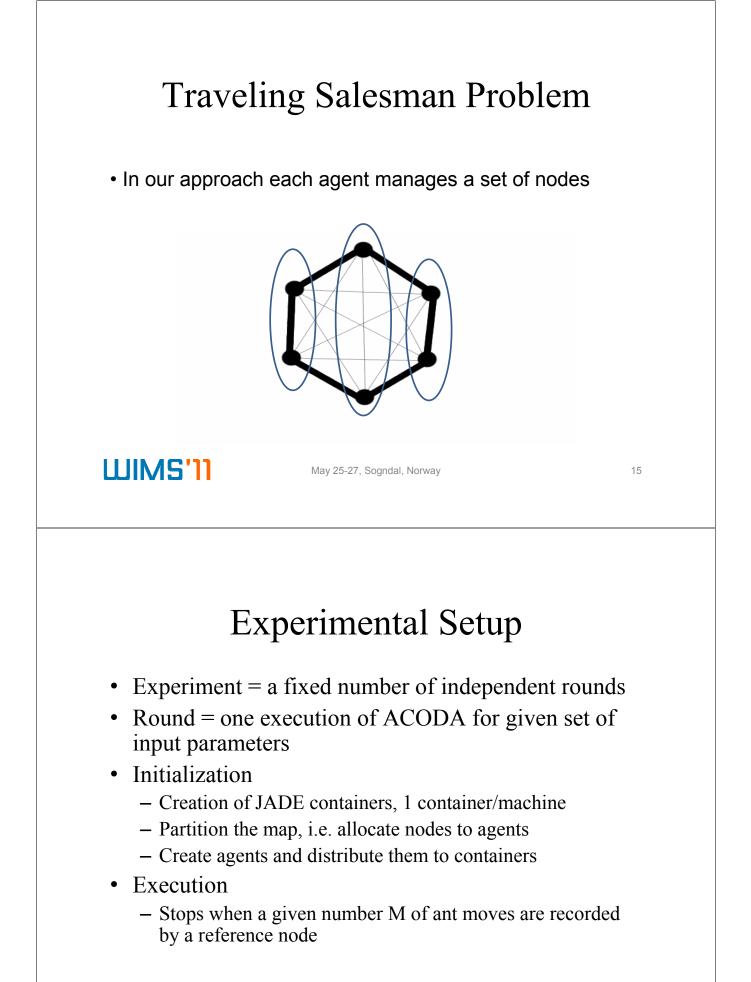
 $\tau_{i,j} = (1 - \zeta) \tau_{i,j} + \zeta \tau_0$

where:

- ζ is the evaporation rate $0 \leq \zeta < 1$
- τ_0 is the initial amount of pheromone on each edge

WIMS'11 May 25-27, Sogndal, Norway 13 Architecture Node Agent receiveAnt() Neighbor -address adjustAttributes() initialize ant -weight -pheromone deposited parameters atanthi Neighbor List at destination set returning read/write defaul Pick a neighbor update tour cost to send ant to







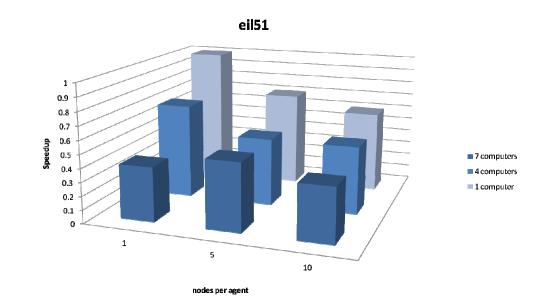
Parameters and Network

- Benchmarks from TSPLIB
 - eil51, st70, kroA100, ch150, gr666
 - $n \in \{51, 70, 100, 150, 666\}$
- $\tau_0 = 1/(n^2 w_{avg})$
- $\rho = \zeta = 0.1, \alpha = 1, \beta = 5$
- *M* = 10000
- Network:
 - 1, 4, and 7 computers with dual core processors at 2.5 GHz and 1GB of RAM memory
 - high-speed Myrinet interconnection network at 2Gb/s
 - variable number of nodes managed by each agent: $k \in \{1, 5, 10\}$

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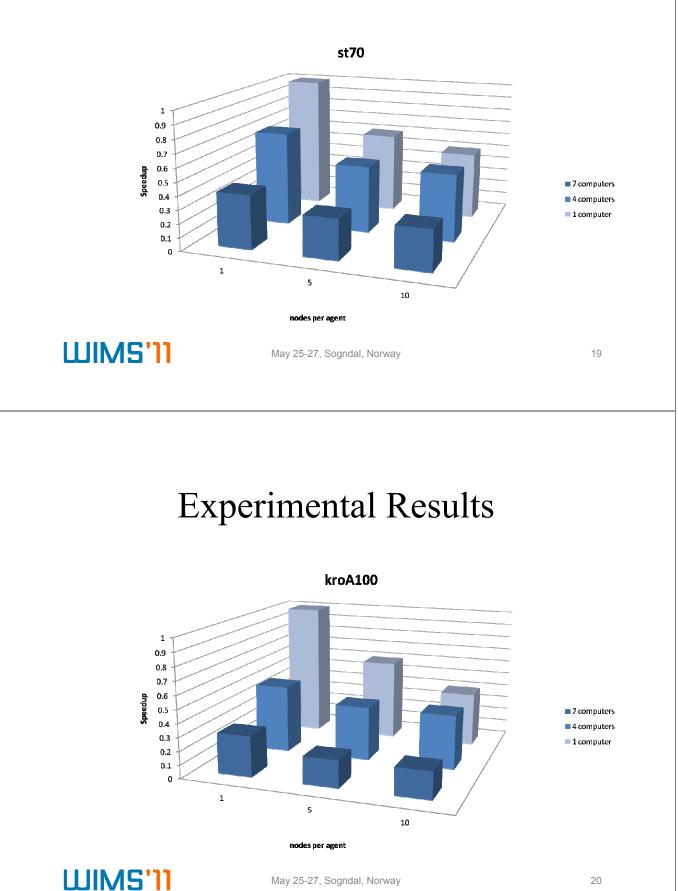
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Experimental Results

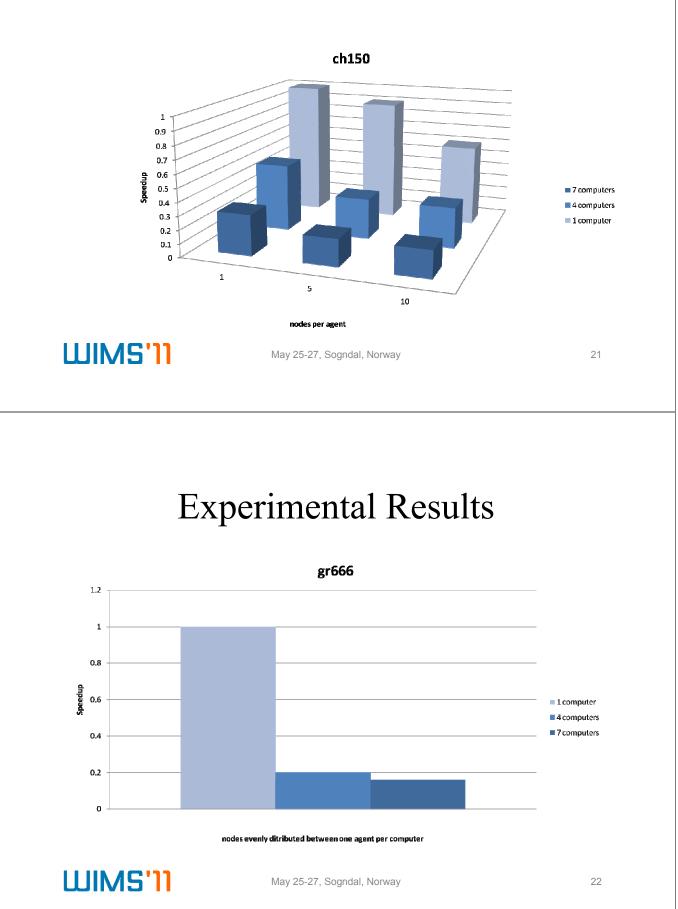




Experimental Results

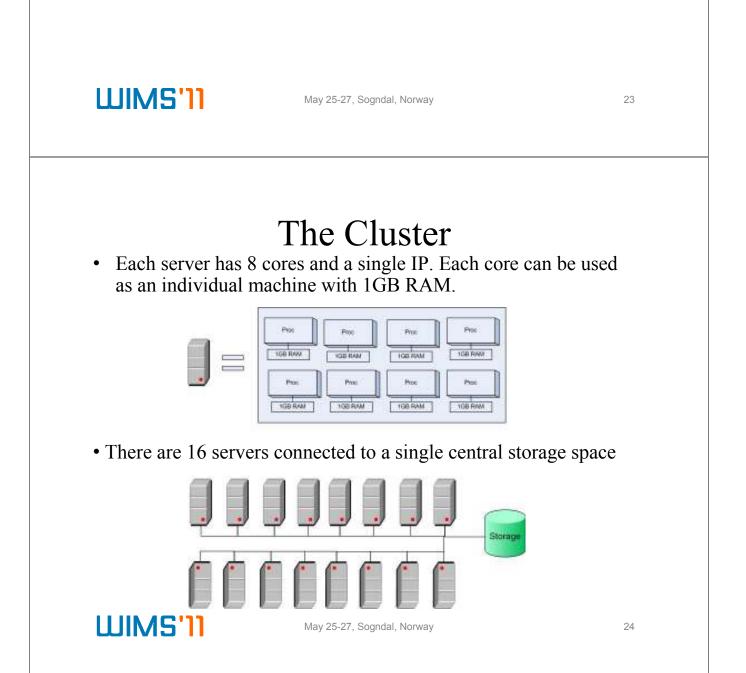


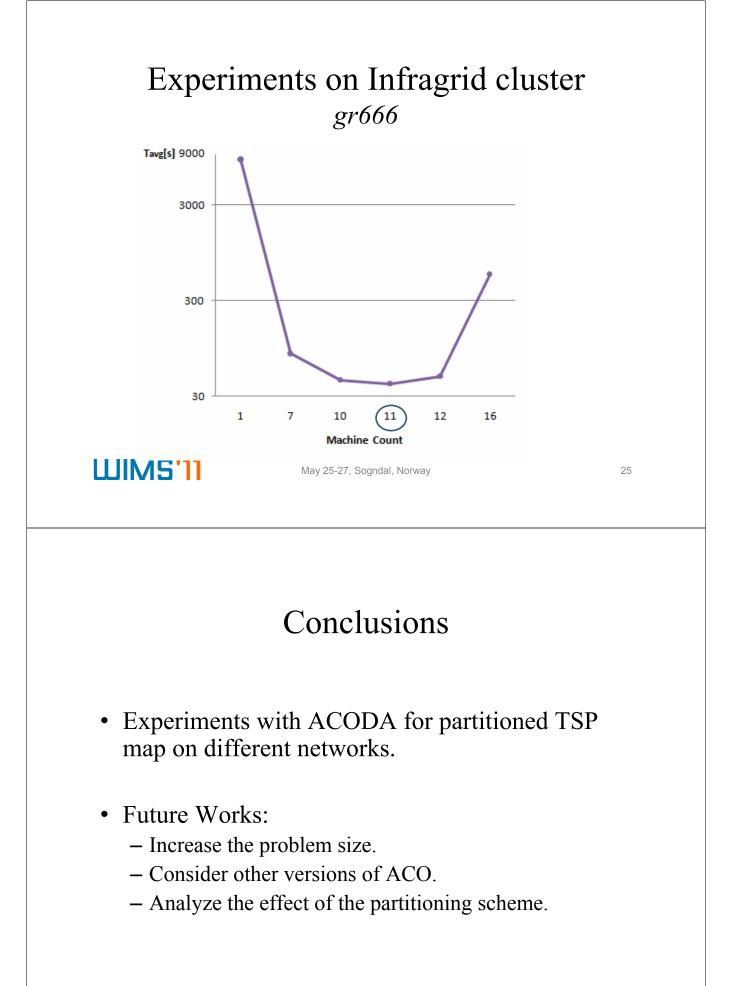
Experimental Results



Recent experiments

• Experiments were also ran on an **Infragrid cluster consisting of 128 cores**, each one with 1GB of RAM, connected by an Infiniband 40 Gb/s network.





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Questions?

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